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Abstract

This article is devoted to the study of monoids which can be endowed with a shuffle product with coefficients in a semiring. We show that, when the multiplicities do not belong to a ring with prime characteristic, such a monoid is a monoid of traces. When the characteristic is prime, we give a decomposition of the congruences $\equiv$ (or relators R) such that $A^*/_{\equiv} = \langle A; R\rangle$ admits a shuffle product. This decomposition involves only addition of primitive elements to the successive quotients. To end with, we study the compatibility with Magnus transformation and examine the case of congruences which are homogeneous for some weight function. The existence of such a weight function is also showed for congruences of depth one.

1 Introduction

Partially commutative structures share many nice combinatorial properties with their free counterpart [8, 19, 21]. They have also been use intensively in the computer science area [12, 13]. These structures can be rapidly described as being presented with

$$< A;\ \{[a,b]\}_{(a,b)\in\theta} >_{\mathbf{cat}}$$

where A is a set of generators (an alphabet), $\theta \in A^2$ an unoriented graph without loop, **cat** a suitable category (Monoid, Group, K-associative algebras, Lie algebras) and $[a,b]$ expresses the fact that a and b commute in the structures of **cat**.

The simplest (and already bearing all the combinatorial power [10]) of these structures is the partially commutative monoid $< A;\ \{ab = ba\}_{(a,b)\in\theta} >_{\mathbf{Mon}}$ (denoted $\mathbb{M}(A,\theta)$) whose algebra over a semiring K ($K < A,\theta >= K\langle\mathbb{M}(A,\theta)\rangle$) is called the algebra of partially commutative polynomials.
It is well known that, if a congruence $\equiv$ (i.e. an equivalence over the free monoid

A^*) is generated by commutations (i.e. is the kernel of a natural morphism $nat : A^* \to \mathbb{M}(A, \theta)$) then the shuffle product of two classes is a class-sum [1] (This can be verified by hand in the general case and in case K is a ring, is related to the fact that $K < A, \theta >$ is the envelopping algebra of the - free - Lie algebra generated by the letters [20]).

This amounts to say that, if two series $S_i = \sum_{w \in A^*} (S_i, w)w;\ i = 1, 2$ are $\equiv$-saturated (i.e. are constant on every class of $\equiv$) then their shuffle product $S_1 \sqcup\!\sqcup S_2$ is again saturated, which in turn reflects, by duality, the fact that there exists a coproduct $c_\equiv : K[A^*/_\equiv] \to K[A^*/_\equiv] \otimes K[A^*/_\equiv]$ such that the square

$$\begin{array}{ccc} K < A > & \xrightarrow{c} & K < A > \otimes K < A > \\ {\scriptstyle nat}\downarrow & & \downarrow{\scriptstyle nat \otimes nat} \\ K[A^*/_\equiv] & \xrightarrow{c_\equiv} & K[A^*/_\equiv] \otimes K[A^*/_\equiv] \end{array} \tag{1}$$

is commutative (c being the coproduct dual to the shuffle). The converse was shown by Duchamp and Krob [7] with the restriction that K be a ring of characteristic 0.

The aim of our paper is to discuss the property defined by (1) which will be called, throughout the paper $K - \sqcup\!\sqcup$ compatibility.
Let us mention here that only the shuffle product is worth and gives rise to such an interesting discussion as the compatibility with the other classical rational laws (sum, external product, star and also Cauchy, Hadamard and infiltration products [6]) over series give weaker or equivalent results (for full details see Commentary 3). The structure of the paper is the following.

First (section 2), we prove that the result of Duchamp and Krob holds in almost every case, the only exception being the rings of prime characteristic for which differences of words that are primitive elements can occur. The iteration of these adjunctions exhausts every finitely generated congruence. More formally, we have:

Theorem 1 *Let $\equiv$ be a finitely generated congruence on A^* which is $K - \sqcup\!\sqcup$ compatible then*

1. *If K is not a ring or if the characteristic of K ($ch(K)$) is not prime then $A^*/_\equiv$ is a partially commutative monoid.*
2. *If $ch(K) = p$ is prime, then a finite $R \subset A^* \times A^*$ exists with a partition $R = \bigcup_{i \in [1,n]} R_i$ such that for $i \in [1, n]$, R_i consists in pairs $(u, v) \in R$ such that (the image of) $u - v$ is primitive in $K[A^*/_{\equiv_{\bigcup_{j=1}^{i-1} R_j}}]$.*

In a second part (section 3), we investigate some properties of the new family of monoids. We prove that these monoids admit a Magnus transformation and we use it to examine the properties of cancellability, gradation, embeddability in a group and roots.

[1] The structure constants are the partially commutative subword coefficients [5, 19].

Theorem 2 *Let $\equiv$ be a finitely generated congruence on A^* which is $K - ш$ compatible and generated by relators of the form $u \equiv v$ where $u - v$ is primitive. We have the following alternatives.*

1. *The monoid $A^*/_\equiv$ is not cancellable.*
2. *A weight function exists $\omega : A \to \mathbb{N}^*$ for which $\equiv$ is homogeneous and $A^*/_\equiv$ embeds in a group.*

2 Compatibility with the Shuffle Product

This section is devoted to study the $K - ш$ compatibility of a congruence. First, we address the problem in the general case (K is a semiring) and give some results common to all the semirings. We treat the case when K is the boolean semiring and when it is a ring, we develop some results about the primitive polynomials which are useful to prove our Theorem 1. In the last paragraph of this section, we sketch the proof.

2.1 General Properties

In the sequel we will denote $\mathbb{N}.1_K$ the subsemiring of a semiring K generated by 1_K.
We first remark that a congruence is $K - ш$ compatible if and only if it is $\mathbb{N}.1_K - ш$ compatible. The property below is straightforward from this observation.

Lemma 3 *Let $\phi : K_1 \to K_2$ be a morphism of semirings then*

1. *If $\equiv$ is $K_1 - ш$ compatible then it is $K_2 - ш$ compatible.*
2. *If ϕ is into the converse holds.*

This shows that, in order to study the $K - ш$ compatibility of a congruence, it suffices to study its $\mathbb{N}.1_K - ш$ compatibility.

Remark 1 *If $\equiv$ is $\mathbb{N} - ш$ compatible then it is $K - ш$ compatible for each semiring K.*

The properties below will be useful in the sequel, their proof are easy and we omit them.

Lemma 4
1. *If $\equiv_1$ and $\equiv_2$ are $K - ш$ compatible congruences then $\equiv_1 \vee \equiv_2$ and $\equiv_1 \wedge \equiv_2$ also are (supremum and infimum are defined with respect to relation "is coarser than").*
2. *Let R be a relator on A^*. The congruence $\equiv_R$ generated by R is $K - ш$ compatible if and only if for each pair $(w_1, w_2) \in R$ we have $c(w_1) \equiv_R^{\otimes 2} c(w_2)$.*

3. *Let $u \in A^+$ and n be the maximal integer such that u can be written as $u = u_1 a^n$ with $a \in A$. Then for each K we have $(c(u), u_1 \otimes a^n) = 1$.*

4. *Each congruence generated by relators of the form $a \equiv b$ (LI) or $dc \equiv cd$ (LC) with $a, b, c, d \in A$ is $K - ш$ compatible.*

5. *Let $B \subseteq A$ be a subalphabet of A. If $\equiv$ is $K - ш$ compatible then its restriction to B^* is.*

2.2 The Boolean Case

This paragraph deals with the case when $K = \mathbb{B}$ (the boolean semiring). The $\mathbb{B} - ш$ compatibility is completely characterized by the following proposition.

Proposition 5 *A congruence is $\mathbb{B} - ш$ compatible if and only if it is generated by relators like $a \equiv 1$ (LE), $a \equiv b$ (LI) or $ab \equiv ba$ (LC), with $a, b \in A$.*

Sketch of the proof. The "if" part is immediate using Properties (1), (2) and (4) of Lemma 4.
For the converse, one defines a suitable section $S \subset A$ such that $S^* \stackrel{nat_S}{\to} A^*/_{\equiv}$ is onto and $\equiv_S = \text{Ker}\ (nat_S)$ is generated by only (LC) relators. We have successively verify that $\equiv_S$ is multihomogeneous and then, in fact, a partially commutative congruence. ■

2.3 Primitive Elements

We suppose now that K is a ring and we examine the link between $K - ш$ compatibility and primitivity of polynomials.

Definition 6 *Let K be a ring and $\equiv$ be a $K - ш$ compatible congruence. A polynomial $P \in K[A^*/_{\equiv}]$ is called primitive if and only if $c_{\equiv}(P) = P \otimes 1 + 1 \otimes P$.*

As usual we have the following property.

Proposition 7 *The submodule $Prim(K[A^*/_{\equiv}])$ of primitive polynomials endowed with the Lie bracket $[\ ,\ ]$ is a Lie algebra.*

We suppose now that $\equiv = \equiv_\theta$ is a relation generated by commutations. Recall first that the free partially commutative monoid is

$$\mathbb{M}(A, \theta) = < A | \{ab = ba\}_{(a,b) \in \theta} > .$$

The monoid $\mathbb{M}(A, \theta)$ can be totally ordered by a relation $<_{std}$ in the following way :

$$t <_{std} t' \Leftrightarrow std(t) <_{lex} std(t')$$

where $std(t)$ denotes the maximal word for the lexicographical order in the commutation class t. Using this order Lalonde [16, 15] has generalized the notion of Lyndon word : the set of Lyndon traces is defined as the set of connected and

primitive[2] traces minimal in their conjugate classes and denoted $Ly(A,\theta)$. In his thesis Lalonde has shown the following theorem.

Theorem *(Lalonde).* *Let $\Lambda : Ly(A,\theta) \to L_K(A,\theta)$ (the - free - Lie algebra generated by the letters in $K < A,\theta >$) be the mapping defined by*

$$\begin{cases} \Lambda(a) = a & \text{if } a \in A \\ \Lambda(l) = [\Lambda(l_1), \Lambda(l_2)] & \text{if } l = l_1 l_2,\ l_1, l_2 \in Ly(A,\theta) \text{ and } |l_2| \text{ minimal} \end{cases}$$

then $(\Lambda(l))_{l \in Ly(A,\theta)}$ is a basis of $L_K(A,\theta)$ as a K module and for each Lyndon trace $l \in Ly(A,\theta)$ one has

$$\Lambda(l) = l + \sum_{t >_{std} l} \beta_t t$$

We have the following proposition.

Proposition 8 *Let p be a prime integer and $u - v$ be a primitive polynomial of $\mathbb{Z}/p\mathbb{Z} < A,\theta >$ such that $u, v \in \mathbb{M}(A,\theta)$. We can prove either $u = a^{p^\alpha}$ and $v = b^{p^\beta}$ or $u = a^{p^\alpha} b^{p^\beta}$ and $v = b^{p^\beta} a^{p^\alpha}$.*

Sketch of the proof. We denote $L^{(p)}_{\mathbb{Z}/p\mathbb{Z}}(A,\theta)$ the Lie algebra of the primitive polynomials. We can prove successively that

a) The only monomials u which are primitive are of the form $u = a^{p^\alpha}$.

b) If $a_1 \cdots a_n - b_1 \cdots b_m$ is a primitive polynomial then we have $a_1 \cdots a_n = a_n \cdots a_1$ and $b_1 \cdots b_m = b_m \cdots b_1$ or $n = m$ and $a_1 \cdots a_n = b_n \cdots {b_1}^3$.

c) Using Lalonde's Theorem, the set $\{(\Lambda(l))^{p^e}\}$ generates $L^{(p)}_{\mathbb{Z}/p\mathbb{Z}}(A,\theta)$ as a K-module and we write $u - v$ in the form

$$u - v = l^{p^e} + \sum_{t > l^{p^e}} \gamma_t t.$$

Without restriction we can consider that $u <_{std} v$ and then $u = l^{p^e}$. According to (b), we have to consider two cases

i) If $u = \overline{u}$ and $v = \overline{v}$, we prove that u and v are primitive and by (a) we get the claim.

ii) Suppose that $u = \overline{v}$. Remarking that, if a power l^e $(e > 0)$ is not of the form a^α or $a^\alpha b^\beta$, one has $l^e = awb^\alpha c^\beta$ with $a, b, c \in A$, $w \in \mathbb{M}(A,\theta)$, $\alpha, \beta > 0$ with $\alpha + \beta$ maximal, $c \neq b$ and $a \neq c$. In this case $(c(l^e), awc^\beta \otimes b^\alpha) = 1$ (by Property (3) of Lemma 4) and $(c(\overline{l^e}), awc^\beta \otimes b^\alpha) = 0$, which proves that if $u \neq a^\alpha b^\beta$ and $u \neq a^\alpha$, $u - v$ is not primitive. Solving the equations $c(a^\alpha - b^\beta) = (a^\alpha - b^\beta) \otimes 1 + (a^\alpha - b^\beta) \otimes 1$ and $c(a^\alpha b^\beta - b^\beta a^\alpha) = (a^\alpha b^\beta - b^\beta a^\alpha) \otimes 1 + 1 \otimes (a^\alpha b^\beta - b^\beta a^\alpha)$, we obtain that α and β are necessarily two powers of p.

[2] In the sense of traces (i.e. a trace is primitive if it can not be written as the power of an other trace)

[3] In the sequel, the mirror image of a trace u will be denoted by $\overline{u}$.

■

2.4 Structure of Compatible Congruences

This paragraph is devoted to sketch the proof of Theorem 1. We recall it here.

Theorem 1. *Let $\equiv$ be a $K - \sqcup\!\sqcup$ compatible and finitely generated congruence on A^*. Then*

1. *If K is not a ring or if the characteristic of K ($ch(K)$) is not prime then $A^*/_{\equiv}$ is a partially commutative monoid.*
2. *If $ch(K) = p$ is prime, then a finite $R \subset A^* \times A^*$ exists with a partition $R = \bigcup_{i\in[1,n]} R_i$ such that for $i \in [1, n]$, R_i consists in pairs (u, v) such that (the image of) $u - v$ is primitive in $K[A^*/_{\equiv_{\cup_{j=1}^{i-1} R_j}}]$.*

Sketch of the proof. We prove first (2). Suppose that A is finite (one can restrict ourselves to the letters of the words of the relators), let K be a ring and $\equiv$ be a finitely generated congruence on A^* which is $K - \sqcup\!\sqcup$ compatible. It is always possible to derive a finite set of relators $R \subset A^* \times A^*$ closed in the following sense

$$u \equiv v \text{ and } \max\{u, v\} < \max\{\max\{u', v'\} | (u', v') \in R\} \Longrightarrow (u, v) \in R.$$

We construct $(R_i)_{i\in[0,n]}$ in the following way.

1. We set $S_0 = R_0 = \emptyset$.
2. For each $i > 0$, R_i is the set of the pairs $(u, v) \in R - \bigcup_{j\leq i-1} S_j$ such that $u - v$ is primitive in $K[A^*/_{\equiv_{\cup_{j\leq i-1} R_j}}]$.
3. The relator S_i is the set of the pairs $(u, v) \in R - \bigcup_{j\leq i-1} R_j$ such that $u \equiv_{R_i} v$.

One can prove that this process ends, remarking that for each i if $\equiv_{\cup_{j\leq i-1} R_j} \neq \equiv_R$ we have $R_i \neq \emptyset$. Now let us prove (1). We may consider two cases.

1. The semiring K is not a ring or a ring of characteristic 0. If $1_K + 1_K = 1_K$, as one has $\mathbb{B} \hookrightarrow K$ Lemma 3 and propositon 5 prove that it is generated by (LE), (LI) or (LC) relators (see Lemma 5). If $1_K + 1_K \neq 1_K$, Lemma 3 implies that $\equiv$ is generated by (LE), (LI) or (LC) relators , examining all these cases, we find that only (LE) is impossible. In the two remaining cases, $A^*/_{\equiv}$ is a free partially commutative monoid.
2. The semiring K is ring of characteristic $n \neq 0$ not prime. We consider two cases.

(a) We have $n \neq p^\alpha$ with p prime and $\alpha > 1$. At least two prime factors p_1 and p_2 of n exist. By Lemma 3 $\equiv$ is $\mathbb{Z}/p_1\mathbb{Z} - \sqcup\!\sqcup$ compatible and $\mathbb{Z}/p_2\mathbb{Z} - \sqcup\!\sqcup$ compatible. Proposition 8 and assertion (2) imply that $\equiv$ is generated only by (LI) or (LC) relators.

(b) We have $n = p^m$, then by Lemma 3 $\equiv$ is $\mathbb{Z}/p^2\mathbb{Z} - \sqcup\!\sqcup$ compatible. Again by Lemma 3 it is $\mathbb{Z}/p\mathbb{Z}$-compatible which implies, using Proposition 8, that the only primitive polynomials $u - v$ are of the form $a^{p^\alpha} - b^{p^\beta}$ or $a^{p^\alpha}b^{p^\beta} - b^{p^\beta}a^{p^\alpha}$. Remarking that $\binom{p^\alpha}{p^{\alpha-1}} \neq 0\ [p^2]$, we find that these relators occur only when $\alpha = \beta = 0$. Using assertion (2), again Proposition 8 gives the result.

■

Such a family $(R_i)_{i\in[1,n]}$ will be called a `primitive partition` of $\equiv$. The minimal length of the primitive partitions will be called the `depth` of $\equiv$.

Example 2 *1. All the congruences generated by relators of the form $a^{p^\alpha} = b^{p^\beta}$ or $a^{p^\alpha}b^{p^\beta} = b^{p^\beta}a^{p^\alpha}$ are $\mathbb{Z}/p\mathbb{Z} - \sqcup\!\sqcup$ compatible with depth 1.*

2. The congruence generated by

$$\begin{cases} a^{2^\alpha}b^{2^\beta}a^{2^\alpha}b^{2^\beta} = b^{2^\beta}a^{2^\alpha}b^{2^\beta}a^{2^\alpha} \\ a^{2^{\alpha+1}}b^{2^\beta} = b^{2^\beta}a^{2^{\alpha+1}} \\ b^{2^{\beta+1}}a^{2^\alpha} = a^{2^\alpha}b^{2^{\beta+1}} \end{cases}$$

with $\alpha, \beta \in \mathbb{N} - \{0\}$ are $\mathbb{Z}/2\mathbb{Z} - \sqcup\!\sqcup$ compatible with depth[4] 2.

3 Group Properties

3.1 Compatibility with Magnus transformation

In this paragraph we show that, when $\equiv$ is a $K - \sqcup\!\sqcup$ compatible congruence (K semiring), one can define a Magnus transformation on $K[A^*/_\equiv]$. Recall that the Magnus transformation is the unique endomorphism of $K < A >$ such that $\mu(a) = 1 + a$ for each letter a. One has here

Lemma 9 *Let K be a semiring and $\equiv$ be a congruence $K - \sqcup\!\sqcup$ compatible. Then it exists an unique morphism $\mu_\equiv$ such that the square*

$$\begin{array}{ccc} K < A > & \xrightarrow{\mu} & K < A > \\ {\scriptstyle nat}\downarrow & & \downarrow{\scriptstyle nat} \\ K[A^*/_\equiv] & \xrightarrow{\mu_\equiv} & K[A^*/_\equiv] \end{array} \qquad (2)$$

is commutative.

[4] We only know congruences with depth 2.

Proof It suffices to remark that $\mu = (Id \otimes ev) \circ c$ where ev is the linear mapping from $K < A >$ on K sending each word $w \in A^*$ on 1. This application is constant over A^*, then it is compatible with the congruence $\equiv$ (i.e. it exists an application $ev_{\equiv}$ sending each class of word to 1). Furthermore, as $\equiv$ is $K - ш$ compatible, $c_{\equiv}$ exists and, therefore our morphism is $\mu_{\equiv} = (Id \otimes ev_{\equiv}) \circ c_{\equiv}$. ■

Remark 3 *Before closing the general case, let us mention that the problem of compatibility can also be addressed for other rational laws. This laws are "." (concatenation or Cauchy product), $*$ (star operation, partially defined),$\times$ (external product) and $+$ (union or sum) for the first kind (which is the realm of Kleene-Schützenberger Theorem [14]) and for the second kind $\odot$ (Hadamard product), ш (shuffle product), $\uparrow$ (infiltration product [18]) and $\uparrow_q$ (q-infiltration product*[5]*, the dual of the coproduct $c_q = (a \otimes 1 + 1 \otimes a) + qa \otimes a$). The results can be summarized as follows*

Kind	Laws	*Compatible with commutation*	*Other*
First	$\times$	*Yes*	*All*
	$+$	*Yes*	*All*
	.	*No*	*?*
	$*$	*No*	*?*
Second	$\odot$	*Yes*	*All*
	ш	*Yes*	*depends of K*
	$\uparrow$	*Yes*	*as* ш
	$\uparrow_q$	*Yes*	*as* ш

3.2 Homogeneous Congruences Compatible with the Shuffle

In this paragraph, we examine the congruences which are homogeneous for some weight function $\omega : A \to \mathbb{N}^+$. The function ω is extended to A^* by $\omega(w) = \sum_{a \in \mathrm{Alph}(w)} |w|_a \omega(a)$, where $|w|_a$ is the partial degree of w with respect to the letter a. Remarking that ω is a morphism the following result is straigthforward.

Lemma 10 *Let $R \subset A^* \times A^*$. The following assertions are equivalent.*

1. *For each $(u, v) \in R$, $\omega(u) = \omega(v)$.*
2. *For each $u, v \in A^*$, $u \equiv_R v \Rightarrow \omega(u) = \omega(v)$.*

Definition 11 *We will say that $\equiv$ is homogeneous (for ω) if and only if it satisfies the assertions of Lemma 10.*

[5] It can be shown that q-infiltration is the only dual law satisfying alphabetical and algebraic constraints [6]

Theorem 2 *Let $\equiv$ be a finitely generated congruence on A^* which is $K - \sqcup\!\sqcup$ compatible and generated by relators of the form $u \equiv v$ where $u - v$ is primitive. We have the following altenative.*

1. *The monoid $A^*/_\equiv$ is not cancellable.*
2. *A weight function exists $\omega : A \to \mathbb{N}^*$ for which $\equiv$ is homogeneous and $A^*/_\equiv$ embeds in a group.*

Sketch of the proof. Suppose that $A^*/_\equiv$ is cancellable. If $ch(K)$ is not prime then $A^*/_\equiv$ is a free partially commutative monoid, and the result is straightforward from [9]. If $ch(K)$ is prime, $\equiv$ is generated by relators of type $a^{p^\alpha} b^{p^\beta} = b^{p^\beta} a^{p^\alpha}$ (pLC) and $a^{p^\alpha} = b^{p^\beta}$ (pLI). As all the (pLC) relators are multihomogeneous it suffices to prove the existence of a weight function for which (pLI) is homogeneous. We prove that, if we have two relators $a^{p^{\alpha_i}} \equiv b^{p^{\beta_i}}$ $(i = 1, 2)$ one has $\alpha_1 - \beta_1 = \alpha_2 - \beta_2$. Denoting by $d_{a,b}$ this difference, the result is a consequence of the following lemma whose proof is easy and left to the reader.

Lemma 12 *Let G be the graph of an equivalence relation on A. Let $d : G \to \mathbb{Z}$ be a function such that*

$$(a, b), (a, c) \in G \Rightarrow d(a, b) + d(b, c) + d(c, a) = 0.$$

Then,

1. *It exists a (potential) function $h : A \to \mathbb{Z}$ such that $d(a, b) = h(b) - h(a)$.*
2. *If A is finite, we can choose h positive.*

End of the proof. We remark that $d_{a,b} + d_{b,c} + d_{c,a} = 0$, according to Lemma 12 we can construct a function $h : A \to \mathbb{N}^+$ such that $d_{a,b} = h(b) - h(a)$. With the weight function $\omega = p^h$, the congruence $\equiv$ is homogeneous. On the other hand we have $\mu_\equiv(A^*) \subset 1 + \mathcal{M}_K(A^*/_\equiv)$ where $\mathcal{M}_K(A^*/_\equiv)$ is the ideal of series S such that $(S, 1) = 0$. But, as $\equiv$ is homogeneous, $1 + \mathcal{M}_K(A^*/_\equiv)$ is a group. Furthermore $\mu_\equiv(w) = w + \sum_{\omega(u)<\omega(w)} n_u u$, which implies that $\mu_\equiv : A^*/_\equiv \hookrightarrow 1 + \mathcal{M}_K(A^*/_\equiv)$ is into. This proves the result. ∎

We now give quickly some properties of the Magnus group $1 + \mathcal{M}_K(A^*/_\equiv)$. It has been shown in [9] that, if $ch(K) = 0$, the function $S \to S^n$; $n > 0$ is one to one within the Magnus group (the monoid then is partially commutative). We cannot expect such a property if $ch(K) = p$ because the congruence $a^p \equiv b^p$ can occur with $a \not\equiv b$. However, we have

Proposition 13 *Let K be a ring of prime characteristic p and $\equiv$ a $K - \sqcup\!\sqcup$ compatible and cancellable congruence. Then:*
i) If $q \not\equiv 0$ $[p]$ the function $S \to S^q$ is one to one within the Magnus group $1 + \mathcal{M}_K(A^*/_\equiv)$.
ii) If $q_i \not\equiv 0$ $[p]$; $i = 1, 2$ then $S^{q_1} T^{q_2} = T^{q_2} S^{q_1}$ implies $ST = TS$

Proof Solving degree by degree

$$(1+\sum_{i=1}^{\infty} X_i)^q = (1+\sum_{i=1}^{\infty} Y_i) \tag{3}$$

(with two auxilliary alphabets such that $deg(X_i) = deg(Y_i)$), it can easily be seen that the series $\sum_{k=0}^{\infty} \binom{\frac{1}{q}}{k} X^k$ has its coefficients in $\mathbb{Z}[1/q]$. This proves (i). Now (ii) is a consequence of (i) as the hypothesis can be reformulated $T^{-q_2} S^{q_1} T^{q_2} = S^{q_1}$.

■

4 Conclusion

We have seen that in "almost every case", the congruences compatible with the shuffle product give partially commutative quotients. The only degenerate case occurs when the characteristic of the ground ring is prime and gives rise to a bunch of new phenomena. A process (primitive partitions) to analyse these new congruences has been described. Moreover, for congruences of depth one, we have a complete description of the relators. This allows to prove, thanks to a "Magnus-type" transformation, that the quotients are either not cancellable or embeddable in a group. This transformation gives us the opportunity to use some analytic tools as the series of the $q-root$ for q prime to the characteristic. An infinite family of congruences of depth two has been provided. The problem of describing higher depth remains however open.

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